

Expander Graphs and Their Applications (V)

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λ , λ_2 and $h(G)$

Recall: $\lambda = \max\{|\lambda_2|, |\lambda_n|\}$

Theorem (**Dodziuk**, 84; **Alon** and **Milman**, 85; **Alon**, 86)

Let G be a d -regular graph with spectrum $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$. Then

$$\frac{d - \lambda_2}{2} \leq h(G) \leq \sqrt{2d(d - \lambda_2)}.$$

Corollary

$(d - \lambda)/2 \leq h(G)$. That is, smaller λ implies bigger expansion.

Does bigger expansion imply smaller λ ?

No! Bipartite graphs. Consequently, we cannot apply, say, Expander Mixing Lemma etc, directly on (combinatorial) expander graphs!

G and \tilde{G}

Let $G = (V, E)$ be a d -regular graph. We define a $(d + 1)$ -regular graph \tilde{G} by *adding a (new) selfloop on every vertex in V* .

- ▶ \tilde{G} is not bipartite.
- ▶ $h(\tilde{G}) = h(G)$.
- ▶ $A(\tilde{G}) = A(G) + I$.
- ▶ The eigenvalues of $A(\tilde{G})$ are $d + 1, \lambda_2 + 1, \dots, \lambda_n + 1$ with the same corresponding eigenvectors v_1, v_2, \dots, v_n .

G and \tilde{G} (cont'd)

Theorem

Let $\{G_i\}_{i \in \mathbb{N}}$ be a family of d -regular expander graphs. Then there exists a constant $\delta < d + 1$ such that $\lambda(\tilde{G}_i) \leq \delta$ for every $i \in \mathbb{N}$.

Proof.

As $h(\tilde{G}_i) = h(G_i) \geq \varepsilon$ for a fixed constant $\varepsilon > 0$, by **Dodziuk, Alon** et.al. we have

$$\lambda_2(\tilde{G}_i) \leq \iota.$$

for a fixed constant $\iota < d + 1$.

If $\lambda(\tilde{G}_i) = \lambda_2(\tilde{G}_i)$, then $\lambda(\tilde{G}_i) \leq \iota$.

If $\lambda(\tilde{G}_i) = -\lambda_n(\tilde{G}_i)$, then $\lambda(\tilde{G}_i) = -(\lambda_n(G_i) + 1) \leq d - 1$.

So we can take $\delta := \max\{\iota, d - 1\}$. □

Bigger expansion of G implies smaller λ of \tilde{G} .

Expander Graphs are Good

Recall:

Corollary

For every $\varepsilon > 0$, (n, d, α) -graphs are (c, d) -good with $c = -3/(2 \log \alpha) + \varepsilon$.

Corollary

Let $\{G_i\}_{i \in \mathbb{N}}$ be a family of d -regular expander graphs. Then G_i s are (c, d) -good for some appropriate $c \in \mathbb{N}$.

Proof.

\tilde{G}_i is a $(n, d + 1, \alpha)$ -graph with a fixed constant $\alpha < 1$.

\tilde{G}_i is $(c, d + 1)$ -good for some fixed $c \in \mathbb{N}$.

G_i is (c, d) -good. □

We can prove the above corollary directly using the definition of $h(G)$.

The Existence of Expander Graphs

Using *probabilistic method* we can prove that

almost every graph is an expander graph.

In particular:

Theorem

For every sufficient large $d \in \mathbb{N}$, there exists a $(d^4, d, 1/4)$ -graph.

The Inductive Construction of Expander Graphs

The (Normal) Graph Product

Let G be a graph with adjacency matrix A and $k \in \mathbb{N}$, then G^2 is a graph with adjacency matrix A^2 . That is, for every pair of vertices u and v

there is an edge between u and v in $G^k \iff$ there is a walk of length k in G .

Theorem

If G is an (n, d, α) -graph, then G^k is an (n, d^k, α^k) -graph.

Inductive Construction Using the Zig-Zag Product

Let G and H be two graphs. We will define their zig-zag product $G \textcircled{Z} H$ in such a way that the following theorem is true.

Theorem (Reingold, Vadhan, and Wigderson, 02)

If G is an (n, m, α) -graph and H an (m, d, β) -graph. Then $G \textcircled{Z} H$ is an $(nm, d^2, \varphi(\alpha, \beta))$ -graph, where the function φ satisfies the following conditions:

- ▶ *If $\alpha < 1$ and $\beta < 1$, then $\varphi(\alpha, \beta) < 1$.*
- ▶ *$\varphi(\alpha, \beta) \leq \alpha + \beta$.*
- ▶ *$\varphi(\alpha, \beta) \leq 1 - (1 - \beta^2)(1 - \alpha)/2$.*

Inductive Construction Using Zig-Zag Products (cont'd)

Let H be a $(d^4, d, 1/4)$ -graph for a constant d . Then we let

$$G_1 := H^2, \quad \text{and} \quad G_{n+1} := G_n^2 \otimes H \text{ for } n \geq 1.$$

Theorem

Every graph G_n is a $(d^{4n}, d^2, 1/2)$ -graph for all $n \in \mathbb{N}$.

Proof.

Trivial for G_1 by the normal graph product.

By induction hypothesis, G_n^2 is a $(d^{4n}, d^4, 1/4)$ -graph. Then by the Zig-Zag Theorem,

$$\lambda(G_{n+1}) \leq \lambda(G_n^2) + \lambda(H) = 1/4 + 1/4 = 1/2.$$



The Definition of the Zig-Zag Product

Let G be an (n, m, α) -graph and H an (m, d, β) -graph. For every vertex $v \in V(G)$ we fix some numbering

$$e_v^1, e_v^2, \dots, e_v^m$$

of the edges incident to v in G . We also assume

$$V(H) = [m] = \{1, 2, \dots, m\}.$$

Definition

$G \textcircled{Z} H = (V(G) \times [m], E')$ where

- there is an edge between (v, i) and (u, j)
- \iff there exist some $k, \ell \in [m]$ such that
 - in H there are edges between i and k, ℓ and j
 - and $e_v^k = e_u^\ell$ in G .

The Adjacency Matrix of the Zig-Zag Product

Every edge between (v, i) and (u, j) in $G \circledast H$ can be viewed a walk length 3 consisting of three steps:

- (i) an edge in the v -th copy in H , i.e., the edge between i and k ;
- (ii) an edge in G , i.e., the edge $e_v^k = e_u^\ell$;
- (iii) an edge in the u -th copy of H , i.e., the edge between ℓ and j .

Let \hat{B} be the normalized adjacency matrix of H . (i) and (iii) are done on n disjoint copies of H with the corresponding transition matrix

$$\tilde{B} := \hat{B} \otimes I_n \quad (\text{tensor product})$$

The Adjacency Matrix of the Zig-Zag Product (cont'd)

In step (ii) we move from a vertex (v, k) to the *unique* vertex (u, ℓ) with $e_v^k = e_u^\ell$. Thus, the corresponding transition matrix P is defined by

$$P_{(v,k),(u,\ell)} := \begin{cases} 1, & \text{if } e_v^k = e_u^\ell \\ 0, & \text{otherwise.} \end{cases}$$

P is a *permutation matrix*, hence for every f

$$\|Pf\|_2 = \|f\|_2.$$

In total, the normalized adjacency matrix of $G \otimes H$ is

$$Z := \tilde{B}P\tilde{B}.$$

A Weak Version of the Zig-Zag Theorem

Theorem

If G is an (n, m, α) -graph and H an (m, d, β) -graph. Then $G \otimes H$ is an (nm, d^2, φ) -graph with

$$\varphi \leq \alpha + \beta + \beta^2.$$

Yet another view of λ

Recall we have proved

$$\lambda_2 = \max \left\{ \frac{\mathbf{v}^T A \mathbf{v}}{\|\mathbf{v}\|_2^2} \mid \mathbf{v} \text{ and } \mathbf{1} \text{ are orthogonal, i.e., } \langle \mathbf{v}, \mathbf{1} \rangle = 0 \right\}.$$

Lemma

$$\lambda = \max \left\{ \frac{|\mathbf{v}^T A \mathbf{v}|}{\|\mathbf{v}\|_2^2} \mid \mathbf{v} \text{ and } \mathbf{1} \text{ are orthogonal, i.e., } \langle \mathbf{v}, \mathbf{1} \rangle = 0 \right\}.$$

Proof.

By $\langle \mathbf{v}, \mathbf{1} \rangle = 0$, $\mathbf{v} = \sum_{2 \leq i \leq n} \beta_i \mathbf{v}_i$. Then

$$\begin{aligned} |\mathbf{v}^T A \mathbf{v}| &= |\langle \mathbf{v}, A \mathbf{v} \rangle| = \left| \left\langle \sum_{2 \leq i \leq n} \beta_i \mathbf{v}_i, \sum_{2 \leq i \leq n} \lambda_i \beta_i \mathbf{v}_i \right\rangle \right| \\ &= \left| \sum_{2 \leq i \leq n} \lambda_i \langle \beta_i \mathbf{v}_i, \beta_i \mathbf{v}_i \rangle \right| \leq \lambda \sum_{2 \leq i \leq n} \langle \beta_i \mathbf{v}_i, \beta_i \mathbf{v}_i \rangle = \lambda \|\mathbf{v}\|_2^2. \end{aligned}$$

The equality is attained when $\mathbf{v} = \mathbf{v}_2$ or $\mathbf{v} = \mathbf{v}_n$. □

Proof

By the previous lemma, it suffices to show that for every mn -vector f with $\langle f, \mathbf{1} \rangle = 0$ we have

$$\frac{|fZf|}{\|f\|_2^2} \leq \alpha + \beta + \beta^2$$

Define f^{\parallel} by

$$f^{\parallel}(x, i) := \frac{\sum_{j \in [m]} f(x, j)}{m}$$

and

$$f^{\perp} := f - f^{\parallel}.$$

It is easy to verify that $\langle f^{\parallel}, f^{\perp} \rangle = 0$, hence

$$\|f\|_2^2 = \|f^{\parallel}\|_2^2 + \|f^{\perp}\|_2^2.$$

Proof (cont'd)

$$\|fZf\| = \left| f\tilde{B}P\tilde{B}f \right| \leq \left| f^{\parallel}\tilde{B}P\tilde{B}f^{\parallel} \right| + 2 \left| f^{\parallel}\tilde{B}P\tilde{B}f^{\perp} \right| + \left| f^{\perp}\tilde{B}P\tilde{B}f^{\perp} \right|$$

- ▶ $\tilde{B}f^{\parallel} = f^{\parallel}$ by $\hat{B}\mathbf{1}_m = \mathbf{1}_m$.
- ▶ $\|\tilde{B}f^{\perp}\|_2 \leq \beta\|f^{\perp}\|_2$ by $\|\hat{B}u\|_2 \leq \beta\|u\|_2$ for every $\langle u, \mathbf{1} \rangle = 0$.

Then

$$\begin{aligned} \|fZf\| &\leq \left| f^{\parallel}Pf^{\parallel} \right| + 2 \left| f^{\parallel}P\tilde{B}f^{\perp} \right| + \left| f^{\perp}\tilde{B}P\tilde{B}f^{\perp} \right| \\ &\leq \left| f^{\parallel}Pf^{\parallel} \right| + 2\beta\|f^{\parallel}\|_2\|f^{\perp}\|_2 + \beta^2\|f^{\perp}\|_2^2. \end{aligned}$$

The second inequality uses the fact that P is a permutation and the fact that $|f_1 f_2| \leq \|f_1\|_2 \|f_2\|_2$.

Proof (cont'd)

Then we show

$$\left| \langle f^\parallel, Pf^\parallel \rangle \right| \leq \alpha \|f^\parallel\|_2^2.$$

Define a function/vector \underline{g} on $V(G)$ by

$$g(v) := \sqrt{m} f^\parallel(v, i) = \frac{\sum_{j \in [m]} f(x, j)}{\sqrt{m}}.$$

By **Cauchy-Schwartz**, $\|g\|_2^2 \leq \|f^\parallel\|_2^2$.

By the definition of P ,

$$f^\parallel Pf^\parallel = g \hat{A} g$$

where \hat{A} is the normalized adjacency matrix of G .

As $\langle f^\parallel, \mathbf{1}_{mn} \rangle = 0$, we have $\langle g, \mathbf{1}_n \rangle = 0$. Therefore, $|g \hat{A} g| \leq \alpha \|g\|_2^2$. The result follows.

Proof (cont'd)

Finally we have

$$\|fZf\| \leq \alpha \|f^{\parallel}\|_2^2 + 2\beta \|f^{\parallel}\|_2 \|f^{\perp}\|_2 + \beta^2 \|f^{\perp}\|_2^2.$$

Recall $\|f\|_2^2 = \|f^{\parallel}\|_2^2 + \|f^{\perp}\|_2^2$,

$$\frac{\|fZf\|}{\|f\|_2^2} \leq \alpha + \beta + \beta^2.$$

□