



# A Introduction to Game Theory

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# Outline

- ✓ 1<sup>st</sup> a brief introduction of Game theory
- ✓ 2<sup>nd</sup> Strategic games
- ✓ 3<sup>rd</sup> Extensive games

# Game theory and CS

- ✓ There has been a remarkable increase in work at the interface of computer science and game theory in the past decade.

----Joseph Y. Halpern

《Computer Science and Game Theory: A Brief Survey》



✓ 1<sup>st</sup> a brief introduction of Game theory

# In the beginning

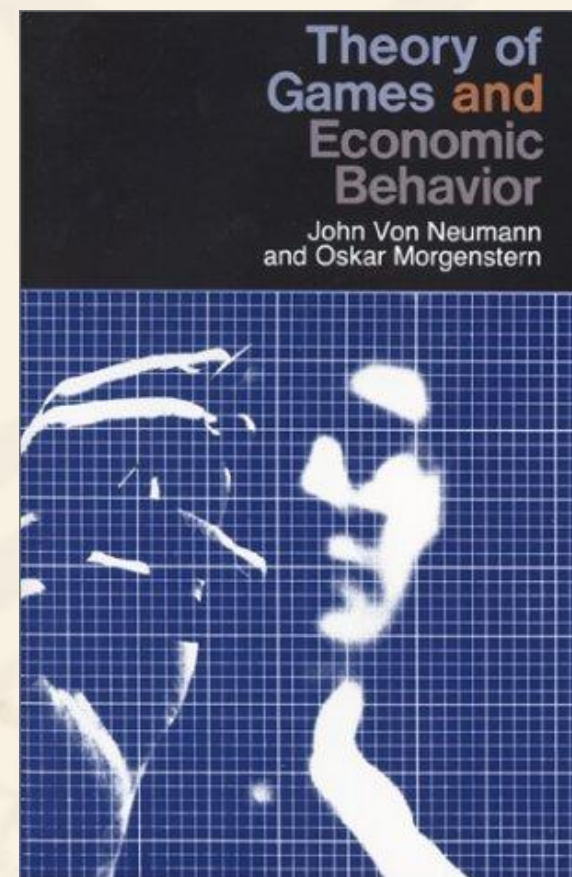


- ✓ Some game-theoretic ideas traced back to the 18-th century.
- ✓ Emile Borel (1871~1956) and John von Neumann (1903~1957) began the major development of game theory.



# The book

- ✓ Game theory became a field since the book of John von Neumann and Oskar Morgenstern published in 1944.



# John Nash(1928- )

- ✓ Received his Ph.D. from Princeton University with a 28-page thesis on his 22-nd birthday.

Invented the notion of Nash equilibrium.

- ✓ Wrote a seminal paper on bargain theory.



# What is Game Theory?

**Game theory is a study of how to mathematically determine the best strategy for given conditions in order to optimize the outcome**



# Rationality

## Assumptions:

- ✓ humans are rational beings
- ✓ humans always seek the best alternative in a set of possible choices

# Why assume rationality?

- ✓ narrow down the range of possibilities
- ✓ predictability

# Types of Games

- ✓ **Perfect vs. Imperfect information**
- ✓ **Sequential vs. Simultaneous moves**
- ✓ **Zero vs. non-zero sum**
- ✓ **Cooperative vs. conflict**

# Types of game and Equilibrium

<div>order</div> <div>Information</div>	Strategic game	Extensive game
perfect	Nash Equilibrium (John Nash)	subgame perfect equilibrium (Reinhard Selten )
imperfect	Bayesian equilibrium (John Harsanyi )	perfect Bayesian equilibrium (Reinhard Selten )



# Two examples

✓ 1) Prisoner's Dilemma

✓ 2) Boxed Pigs

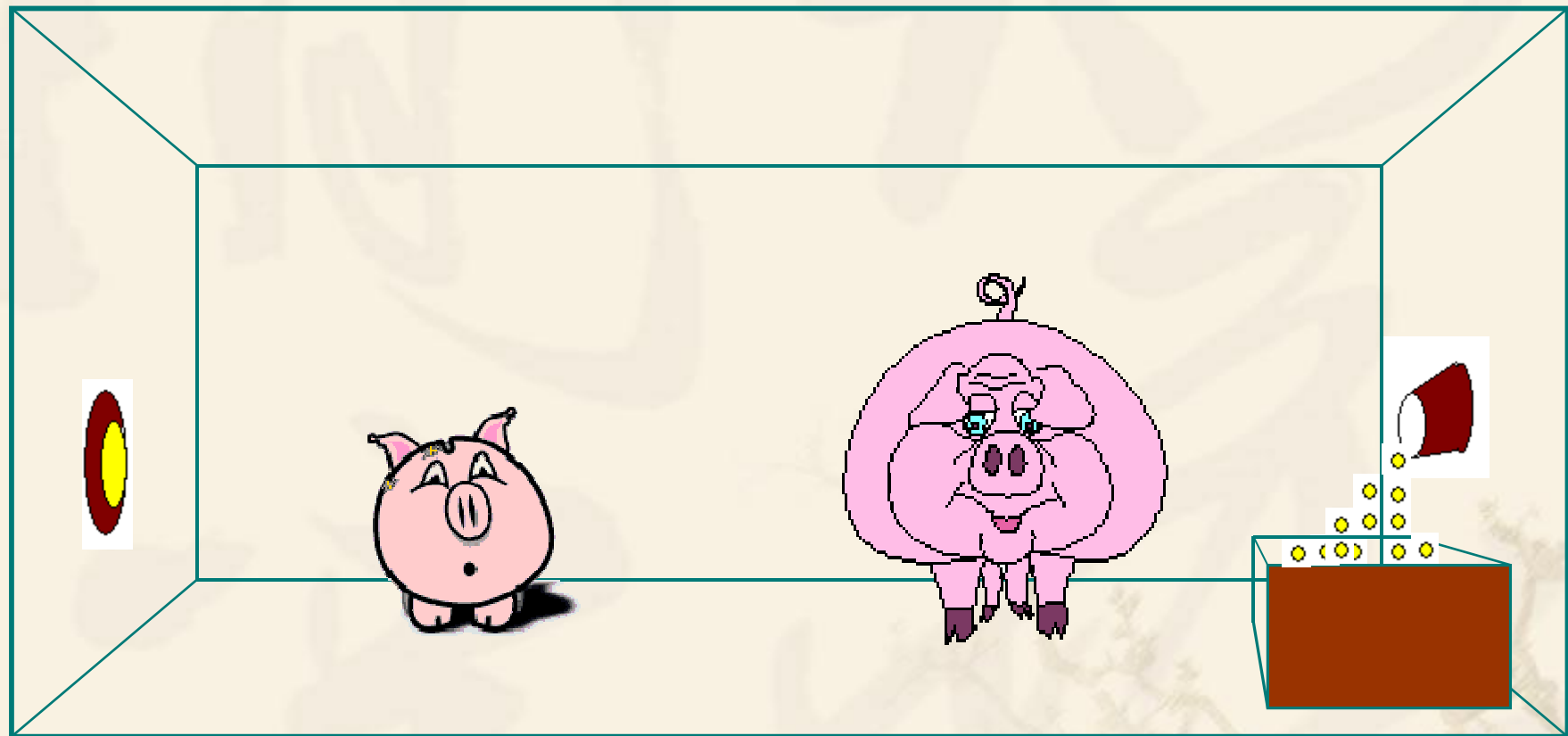
# Prisoner's Dilemma



# Prisoner's Dilemma

		Prisoner 2	
		<i>Blame</i>	<i>Don't</i>
Prisoner 1	Blame	10 , 10	0 , 20
	Don't	20 , 0	1 , 1

# Boxed Pigs



Cost to press  
button = 2 units

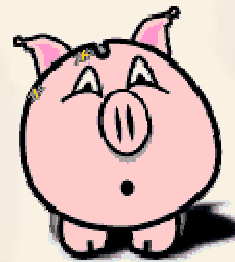
When button is pressed,  
food given = 10 units



# Decisions, decisions...



Big  
Pig



Little Pig

**Press**

**Wait**

***Press***  
**5 , 1**

**9 , -1**

***Wait***  
**4 , 4**

**0 , 0**

## 2<sup>nd</sup> Strategic game

$v < N, (A_i), (u_i) >$

A **strategic game** consists of

- a finite set  $N$  (the set of **players**)
- for each player  $i \in N$  a nonempty set  $A_i$  (the set of **actions** available to player  $i$ )
- for each player  $i \in N$  a preference relation  $\succsim_i$  on  $A = \times_{j \in N} A_j$  (the **preference relation** of player  $i$ ).

If the set  $A_i$  of actions of every player  $i$  is finite then the game is *finite*.

# Nash Equilibrium

**“If there is a set of strategies with the property that no player can benefit by changing her strategy while the other players keep their strategies unchanged, then that set of strategies and the corresponding payoffs constitute the Nash Equilibrium. ”**

# Nash Equilibrium

A Nash equilibrium of a strategic game  $\langle N, (A_i), (\succeq_i) \rangle$ , is a profile  $\mathbf{a}^* \in A$  of actions with the property that for every player  $i \in N$  we have

$$(a_{-i}^*, a_i^*) \succeq_i (a_{-i}^*, a_i) \text{ for all } a_i \in A_i.$$



# Imperfect Information

- ✓ **Partial or no information concerning the opponent is given in advance to the player's decision.**

# Bayesian game

A **Bayesian game** consists of  $\langle N, \Omega, (A_i), (T_i), (\tau_i), (p_i), (\succeq_i) \rangle$

- a finite set  $N$  (the set of **players**)

- a finite set  $\Omega$  (the set of **states**)

and for each player  $i \in N$

- a set  $A_i$  (the set of **actions** available to player  $i$ )

- a finite set  $T_i$  (the set of **signals** that may be observed by player  $i$ ) and a function  $\tau_i: \Omega \rightarrow T_i$  (the **signal function** of player  $i$ )

- a probability measure  $p_i$  on  $\Omega$  (the **prior belief** of player  $i$ ) for which  $p_i(\tau_i^{-1}(t_i)) > 0$  for all  $t_i \in T_i$

- a preference relation  $\succeq_i$  on the set of probability measures over  $A \times \Omega$  (the **preference relation** of player  $i$ ), where  $A = \times_{j \in N} A_j$ .

# BoS

Two people wish to go out together to a concert of music by either Bach or Stravinsky. Their main concern is to go out together, but one person prefers Bach and the other person prefers Stravinsky. Representing the individuals' preferences by payoff functions

	<i>Bach</i>	<i>Stravinsky</i>
<i>Bach</i>	2, 1	0, 0
<i>Stravinsky</i>	0, 0	1, 2

## Variant of *BoS* with imperfect information

Consider another variant of the situation modeled by *BoS*, in which neither player knows whether the other wants to go out with her.

Specifically, suppose that player 1 thinks that with probability  $\frac{1}{2}$  player 2 wants to go out with her, and with probability  $\frac{1}{2}$  player 2 wants to avoid her, and player 2 thinks that with probability  $\frac{2}{3}$  player 1 wants to go out with her and with probability  $\frac{1}{3}$  player 1 wants to avoid her.

Assume that each player knows her own preferences.



## Variant of *BoS* with imperfect information

*Players* The pair of people.

*States* The set of states is  $\{yy, yn, ny, nn\}$ .

*Actions* The set of actions of each player is  $\{B, S\}$ .

*Signals* Player 1 receives one of two signals,  $y_1$  and  $n_1$ ; her signal function  $\tau_1$  satisfies  $\tau_1(yy) = \tau_1(yn) = y_1$  and  $\tau_1(ny) = \tau_1(nn) = n_1$ . Player 2 receives one of two signals,  $y_2$  and  $n_2$ ; her signal function  $\tau_2$  satisfies  $\tau_2(yy) = \tau_2(ny) = y_2$  and  $\tau_2(yn) = \tau_2(nn) = n_2$ .



## Variant of *BoS* with imperfect information

*Beliefs* Player 1 assigns probability  $\frac{1}{2}$  to each of the states  $yy$  and  $yn$  after receiving the signal  $y_1$  and probability  $\frac{1}{2}$  to each of the states  $ny$  and  $nn$  after receiving the signal  $n_1$ . Player 2 assigns probability  $\frac{2}{3}$  to the state  $yy$  and probability  $\frac{1}{3}$  to the state  $ny$  after receiving the signal  $y_2$ , and probability  $\frac{2}{3}$  to the state  $yn$  and probability  $\frac{1}{3}$  to the state  $nn$  after receiving the signal  $n_2$ .

# Payoffs

$\frac{1}{2}$

$\frac{1}{2}$

$\frac{2}{3}$

	B	S
B	2, 1	0, 0
S	0, 0	1, 2

State  $yy$

1:  $y_1$

$\frac{2}{3}$

	B	S
B	2, 0	0, 2
S	0, 1	1, 0

State  $yn$

2:  $y_2$

$\frac{1}{2}$

1:  $n_1$

$\frac{1}{2}$

2:  $n_2$

$\frac{1}{3}$

	B	S
B	0, 1	2, 0
S	1, 0	0, 2

State  $ny$

1:  $n_1$

$\frac{1}{3}$

	B	S
B	0, 0	2, 2
S	1, 1	0, 0

State  $nn$

1:  $n_1$

$\frac{1}{2}$

# Bayesian Equilibrium

A Nash equilibrium of a Bayesian game  $(N, \Omega, (A_i), (T_i), (\tau_i), (p_i), (\zeta_i))$  is a Nash equilibrium of the strategic game defined as follows.

- The set of players is the set of all pairs  $(i, t_i)$  for  $i \in N$  and  $t_i \in T_i$ .
- The set of actions of each player  $(i, t_i)$  is  $A_i$ .
- The preference ordering  $\succsim_{(i, t_i)}^*$  of each player  $(i, t_i)$  is defined by

$$a^* \succsim_{(i, t_i)}^* b^* \text{ if and only if } L_i(a^*, t_i) \succsim_i L_i(b^*, t_i),$$

where  $L_i(a^*, t_i)$  is the lottery over  $A \times \Omega$  that assigns probability  $p_i(\omega)/p_i(\tau_i^{-1}(t_i))$  to  $((a^*(j, \tau_j(\omega)))_{j \in N}, \omega)$  if  $\omega \in \tau_i^{-1}(t_i)$  zero otherwise.

# Bayesian Equilibrium

- ✓ (B, y) (B, y)
- ✓ (B, y) (S, n)
- ✓ (B, n) (B, y)
- ✓ (B, n) (S, n)



v3<sup>rd</sup>

# Extensive games



# Extensive game with perfect information

An **extensive game with perfect information** has the following components.

- A Set  $N$  (the set of **players**).
- A set  $H$  of sequences (finite or infinite) that satisfies the following three properties
- The empty sequence  $\emptyset$  is a member of  $H$ .
- $(a^k)_{k=1,\dots,K} \in H$  (where  $K$  may be infinite) and  $L < K$  then  $(a^k)_{k=1,\dots,L} \in H$
- If an infinite sequence  $(a^k)_{k=1}^\infty$  satisfies  $(a^k)_{k=1,\dots,L} \in H$  satisfies  $(a^k)_{k=1,\dots,L} \in H$  for every positive integer  $L$  then  $(a^k)_{k=1}^\infty \in H$ .

# An example

Two people use the following procedure to share two desirable identical indivisible objects. One of them proposes an allocation, which the other then either accepts or rejects. In the event of rejection, neither person receives either of the objects. Each person cares only about the number of objects he obtains.

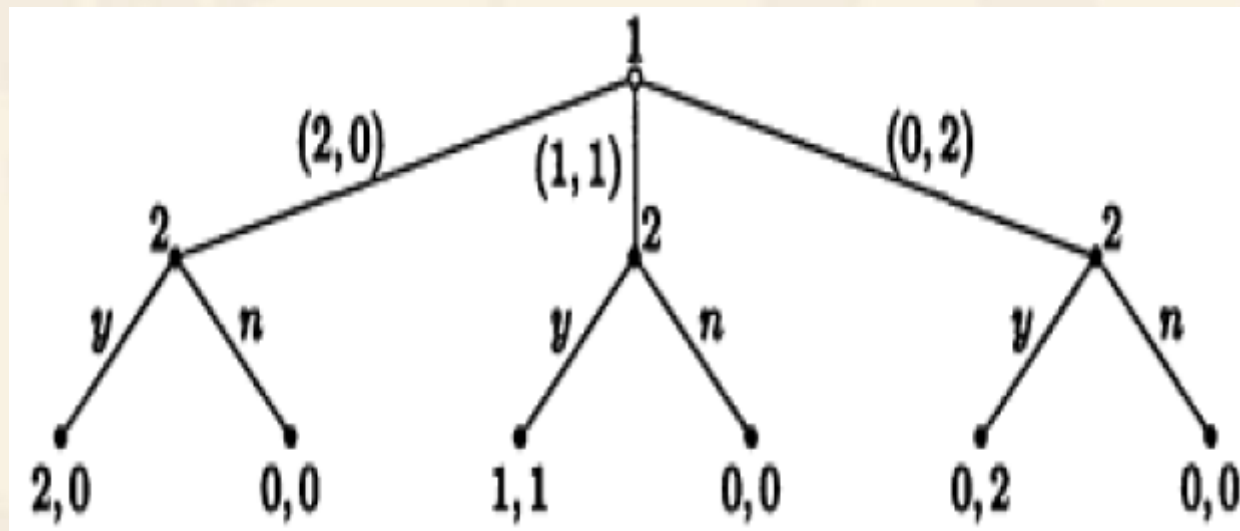
# allocating two identical indivisible objects

An extensive game that models the individuals' predicament is  $\langle N, H, P, (\succeq_i) \rangle$  where

- $N = \{1, 2\}$ ;
- $H$  consists of the ten histories  $\emptyset, (2, 0), (1, 1), (0, 2), ((2, 0), y), ((2, 0), n), ((1, 1), y), ((1, 1), n), ((0, 2), y), ((0, 2), n)$ ;
- $P(\emptyset) = 2$  and  $P(h) = 1$  for every nonterminal history  $h : h \neq \emptyset$ .

$$\begin{aligned} & ((2, 0), y) \succeq_1 ((1, 1), y) \succeq_1 ((0, 2), y) \sim_1 ((2, 0), n) \sim_1 ((1, 1), n) \sim_1 \\ & ((0, 2), n) \text{ and } ((0, 2), y) \succeq_2 ((1, 1), y) \succeq_2 ((2, 0), y) \sim_2 ((0, 2), n) \sim_2 \\ & ((1, 1), n) \sim_2 ((2, 0), n). \end{aligned}$$

# Payoff tree





# Strategy

A strategy of player  $i \in N$  in an extensive game with perfect information  $\langle N, H, P, (z_i) \rangle$  is a function that assigns an action in  $A(h)$  to each nonterminal history  $h \in H \setminus Z$  for which  $P(h) = i$ .

In the just example:

The strategies of player 1:  $\{(2,0), (1,1), (0,2)\}$

The strategies of player 2:  $\{yyy, yyn, yny, ynn, nyy, nyn, nny, nnn\}$



# Nash Equilibrium

For each strategy profile  $s = (s_i)_{i \in N}$  in the extensive game  $(N, H, P, (\succeq_i))$  we define the **outcome**  $O(s)$  of  $s$  to be the terminal history that results when each player  $i \in N$  follows the precepts of  $s$ . That is,  $O(s)$  is the (possibly infinite) history  $(a^1, \dots, a^K) \in Z$  such that for  $0 \leq k < K$  we have  $s_{P(a^1, \dots, a^k)}(a^1, \dots, a^k) = a^{k+1}$ .

A **Nash equilibrium** of an extensive game with perfect information  $(N, H, P, (\succeq_i))$  is a strategy profile  $s^*$  such that for every player  $i \in N$  we have

$$O(s_{-i}^*, s_i^*) \succeq_i O(s_{-i}^*, s_i) \text{ for every strategy } s_i \text{ of player } i.$$

# a problem

- ✓ The Nash equilibria of the example game are  $((2, 0), yyy)$ ,  $((2, 0), yyn)$ ,  $((2, 0), yny)$ ,  $((2, 0), ynn)$ ,  $((1, 1), nyy)$ ,  $((1, 1), nyn)$ ,  $((0, 2), nny)$ ,  $((2, 0), nny)$ , and  $((2, 0), nnn)$ .

# Subgame extensive game

The subgame of the extensive game with perfect information  $\Gamma = (N, H, P, (\xi_i))$  that follows the history  $h$  is the extensive game  $\Gamma(h) = (N, H|_h, P|_h, (\xi_i|_h))$ , where  $H|_h$  is the set of sequences  $h'$  of actions for which  $(h, h') \in H$ , is defined by  $h' \in H|_h$  for each  $\xi_i|_h$  and  $h' \xi_i|_h h''$  is defined by  $(h, h') \xi_i (h, h'')$  if and only if  $h' \in H|_h$ .

# Subgame Perfect Equilibrium

A subgame perfect equilibrium of an extensive game with perfect information  $i \in N$  is a strategy profile  $s^*$  such that for every player  $i \in N$  and every nonterminal history  $h \in H \setminus Z$  for which  $P(h) = i$  we have

$$O_h(s_{-i}^*|h, s_i^*|h) \succeq_i |h| O_h(s_{-i}^*|h, s_i)$$

for every strategy  $s_i$  of player  $i$  in the subgame  $F(h)$ .



# Extensive game with imperfect information

$\langle N, H, P, f_c, (\mathcal{I}_i), (z_i) \rangle$

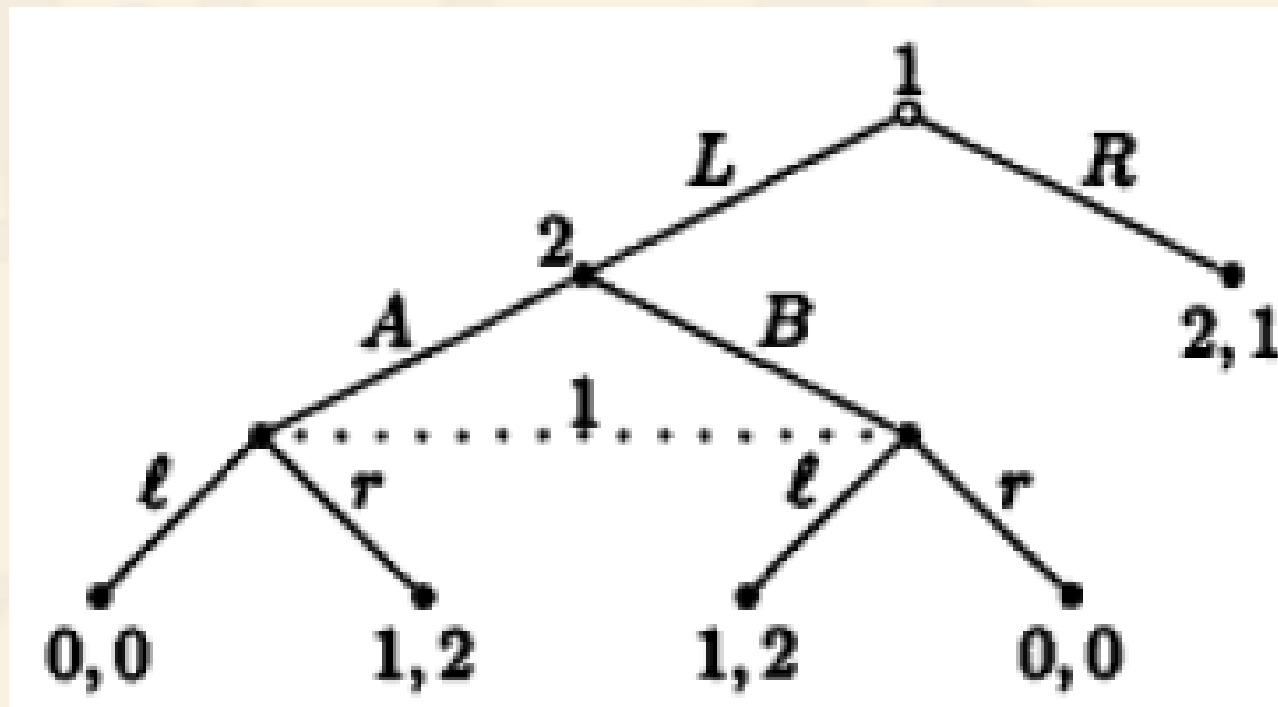
- A function  $P$  that assigns to each nonterminal history (each member of  $H \setminus Z$ ) a member of  $N \cup \{c\}$ . ( $P$  is the **player function**,  $P(h)$  being the player who takes an action after the history  $h$ . If  $P(h) = c$  then chance determines the action taken after the history  $h$ .)
- A function  $f_c$  that associates with every history  $h$  for which  $P(h) = c$  a probability measure  $f_c(\cdot|h)$  on  $A(h)$ , where each such probability measure is independent of every other such measure. ( $f_c(a|h)$  is the probability that  $a$  occurs after the history  $h$ .)



# Extensive game with imperfect information

- For each player  $i \in N$  a partition  $\mathcal{I}_i$  of  $\{h \in H: P(h) = i\}$  with the property that  $A(h) = A(h')$  whenever  $h$  and  $h'$  are in the same member of the partition. For  $I_i \in \mathcal{I}_i$  we denote by  $A(I_i)$  the set  $A(h)$  and by  $P(I_i)$  the player  $P(h)$  for any  $h \in I_i$ . ( $\mathcal{I}_i$  is the **information partition** of player  $i$ ; a set  $I_i \in \mathcal{I}_i$  is an **information set** of player  $i$ .)

# An example



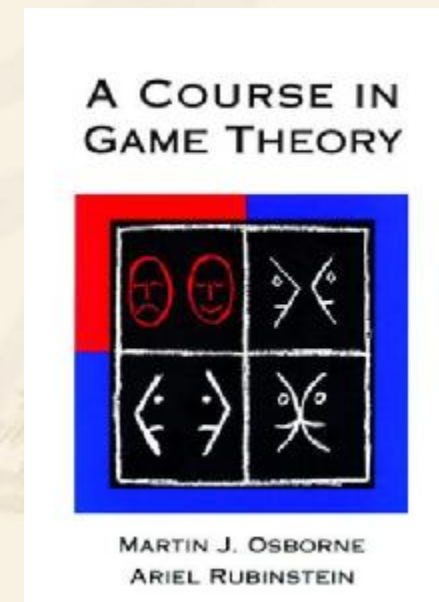
# Nash equilibrium

A Nash equilibrium in mixed strategies of an extensive game is (as before) a profile  $\sigma^*$  of mixed strategies with the property that for every player  $i \in N$  we have

$U(\sigma_{-i}^*, \sigma_i^*) \succeq_i U(\sigma_{-i}^*, \sigma_i)$  for every mixed strategy  $\sigma_i$  of player  $i$ .

# Four parts of the textbook

- ✓ Strategic games (Part I)
- ✓ Extensive games with perfect information (Part II)
- ✓ Extensive games without perfect information (Part III)
- ✓ Coalitional games (Part IV)







Thank you!